Phase Transition of Inhomogeneous Random Graphs

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Joint work with Vlady Ravelomanana

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- Asymptotic analysis of the number of inhomogeneous graphs (some differences with the original model) following the approach of the giant paper [Janson Knuth Łuczak Pittel 93]
- Phase transition of the modeled problems probability for a graph to be bipartite [Pittel Yeum 10], probability of satisfiability of a quantified 2-XOR-SAT formula [Creignou Daudé Egly 07]



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• Each vertex v receives a color c(v),



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- the edges are weighted according to the color of their ends using R = (⁰₁, ¹₀),
- weight $\frac{1}{2}$ for each connected component.

weight(
$$c(G)$$
) := $\left(\frac{1}{2}\right)^{cc(G)} \prod_{(a,b)\in E(G)} R_{c(a),c(b)}$



weight = 0



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weight = $\frac{1}{4}$

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weight(c(G)) :=
$$\left(\frac{1}{2}\right)^{cc(G)} \prod_{(a,b)\in E(G)} R_{c(a),c(b)}$$

$$\sum_{c} weight(c(G)) = \begin{cases} 1 & \text{if } G \text{ is bipartite,} \\ 0 & \text{otherwise.} \end{cases}$$

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The number of (n, m)-bipartite graphs is

$$g_{\begin{pmatrix} 0 & 1\\ 1 & 0 \end{pmatrix}, \frac{1}{2}}(n, m) := \sum_{(n, m) \text{-graph } G} \sum_{c} \text{weight}(c(G)).$$

 $R \in \operatorname{Sym}_{q \times q}(\mathbb{R}_{\geq 0})$ and $\sigma > 0$. A (R, σ) -graph is:

- a vertex colored graph c(G),
- with weight $R_{c(s),c(t)}$ on each edge (s,t),
- $\bullet\,$ and weight σ for each connected component.

weight(
$$c(G)$$
) := $\sigma^{cc(G)} \prod_{(a,b)\in E(G)} R_{c(a),c(b)},$
$$g_{R,\sigma}(n,m) := \sum_{(n,m)\text{-graph } G} \sum_{c} \text{weight}(c(G)).$$

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weight = $\sigma^2 R_{1,1} R_{1,2}^3 R_{1,3}^2 R_{2,3}^2$

Quantified 2-Xor-Sat Formulas [Creignou Daudé Egly 07]

$$\forall x, y, \exists a, b, \dots, h, a \oplus b = x, a \oplus h = y, a \oplus c = x,$$
$$b \oplus e = x, d \oplus f = x, d \oplus g = y, e \oplus h = y$$

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d 1,1

g 1,0

0.0

0,0

с 1.0 ь

1.0

h 0,1



satisfiable iff each cycle contains an even number of *x* and *y*.



The number of satisfiable quantified 2-Xor-Sat formulas with *n* existantial variables and *m* clauses is $g_{R,\sigma}(n,m)$.

Sub-Critical Density of Edges

When $\frac{m}{n} < c(1 - \epsilon)$ and $n \to \infty$, with high probability a (n, m)- (R, σ) -graph consists of trees and unicycle components.



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rooted tree $T_i(z) = z \exp(\vec{R}_i \vec{T}(z))$

Symbolic method

$$z\partial \vec{T} = (I - \operatorname{diag}(\vec{T})R)^{-1}\vec{T}$$

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$$T_i(z) = z \exp(\overrightarrow{R}_i \overrightarrow{T}(z))$$
 $\overrightarrow{T} \sim \overrightarrow{t_0} - \overrightarrow{t_1} \sqrt{1 - \frac{z}{\rho}} + \dots$

Drmota-Lalley-Wood Theorem

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unrooted tree $U = \overleftarrow{1} \overrightarrow{T} - \frac{1}{2} \overleftarrow{T} \overrightarrow{R} \overrightarrow{T}$ $\sim u_0 + u_2(1 - \frac{z}{\rho}) + u_3(1 - \frac{z}{\rho})^{3/2}$

Dissymmetry Theorem

 $z \partial U = T_1 + \ldots + T_q$

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unrooted tree
$$U = \vec{1}\vec{T} - \frac{1}{2}\vec{T}R\vec{T} \qquad \sim u_{0} + u_{2}(1 - \frac{z}{\rho}) + u_{3}(1 - \frac{z}{\rho})^{3/2}$$

unicycle
component
$$V = -\frac{1}{2}\log(\det(I - \operatorname{diag}(\vec{T})R))$$

linear algebra

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$$g_{R,\sigma}(n,m) \sim n! [z^n] \frac{(\sigma U)^{n-m}}{(n-m)!} e^{\sigma V}$$

Large Power scheme [Flajolet Sedgewick 09]: one dominant saddle point.

When $\frac{m}{n} = c(1 + \mu n^{-1/3})$ where $\mu = O(1)$, with high probability a (n, m)- (R, σ) -graph consists of

- trees and unicycle components,
- a cubic multigraph where the vertices are replaced by rooted trees and the edges by paths of trees.

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In a cubic graph, each vertex owns 3/2 edges.

For the ordinary graphs, the gf of the developped cubic part is

$$\operatorname{GF}_{\operatorname{cubic}}\left(z\leftarrow T(z)\left(\frac{1}{1-T(z)}\right)^{3/2}\right)$$

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TreePath_{i,j} =
$$\left(R(I - \operatorname{diag}(\vec{T})R)^{-1}\right)_{i,j}$$

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$$g_{R,\sigma}(n,m) \sim n! [z^n] \sum_{k} \frac{(\sigma U)^{k+n-m}}{(k+n-m)!} \exp(\sigma V) \frac{e_k^{(\sigma)} (T_1 p_1^3 + \dots + T_q p_q^3)^{2k}}{\det(I - \operatorname{diag}(\vec{T})R)^{3k}}$$

Coallescence of two saddle points at the dominant singularity of T(z) [Janson Knuth Łuczak Pittel 93] or [Banderier Flajolet Schaeffer Soria 01].

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Results

New proof of the result of Pittel and Yeum on the probability of bipartiteness, new result on the probability of satisfiability of quantified 2-Xor-Sat formulas:

$$\frac{2m}{n} < 1 - \varepsilon : \mathbb{P}(\operatorname{sat}) \sim \frac{\left(1 - \frac{2m}{n}\right)^{3/8}}{\left(1 + \frac{2m}{n}\right)^{1/8}} \sqrt{1 - \frac{m}{n}},$$
$$\frac{2m}{n} = 1 + \mu n^{-1/3} : \mathbb{P}(\operatorname{sat}) \sim \frac{\Phi_{1/4}(\mu)}{(2n)^{1/8}}.$$

$$\Phi_{\sigma}(\mu) = \sqrt{2\pi} \sum_{k} \frac{e_{k}^{(\sigma)}}{4^{k}} A(3k + \sigma/2, \mu)$$
$$e_{k}^{(\sigma)} = [z^{2k}] \left(\sum_{n} \frac{(6n)! z^{2n}}{(2n)! (3n)! 2^{n} (3!)^{n}} \right)^{\sigma}$$
$$A(y, \mu) = \frac{e^{-\mu^{3}/6}}{3^{(y+1)/3}} \sum_{k} \frac{(3^{2/3} \mu/2)^{k}}{k! \Gamma((y+1-2k)/3)}$$

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- applications to the analysis of algorithms,
- generalisation to hypergraphs.

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